

# Database Security

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## Outline

- 1 Introduction
- 2 Discretionary Access Control  
Granting and Revoking Privileges  
Trojan Horse Attack
- 3 Mandatory Access Control  
The Bell-LaPadula Model  
Multilevel Relations
- 4 Summary

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## Related Concepts

**Authentication:** confirming the identity of users (or programs)

**Authorization\*:** specifying access rights to resources

**Encryption:** encoding data to prevent unauthorized persons  
from reading it (if they managed to access it)

\*Our topic today.

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## Objectives in Securing a Database System

- Secrecy:** protection of data against unauthorized disclosure
- e.g. a student cannot see other students' grades

- Integrity:** prevention of unauthorized data modification
- e.g. only the instructor may assign grades

- Availability:** ensuring authorized access is possible
- e.g. students are not denied seeing their own grades

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## Access Control in a Database System

A **security policy** specifies who is authorized to do what in the system.

- A DBMS provides **access control** mechanisms to help implement a security policy.
- Two complementary types of mechanisms:
  - ① *Discretionary access control*
  - ② *Mandatory access control*

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## Discretionary Access Control

**Idea**

Achieve security based on:

- ① *privileges (certain access rights for tables, columns, etc.), and*
- ② *a mechanism for granting and revoking such privileges at a user's own discretion*

**Authorization administration policy:** specifies how granting/revoking of privileges is organized (i.e. who may grant and revoke)

- *Centralized administration:* only some privileged users
- *Ownership-based administration:* creator of the object

**Administration delegation:** If authorized to do so, a user may assign other users the right to grant or revoke.

In SQL-92, privileges are given to users.  
In SQL:1999, privileges are given to *roles*; those are assigned to users.

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## Granting and Revoking Privileges in SQL

GRANT privileges ON object TO users [WITH GRANT OPTION]

- Possible privileges:
  - SELECT
  - INSERT (column)
  - UPDATE (column)
  - DELETE
  - REFERENCES (column)
- WITH GRANT OPTION allows user to pass on privilege (with or without passing on grant option)

REVOKE [GRANT OPTION FOR] privileges ON object  
FROM users { RESTRICT | CASCADE }

- When a privilege is revoked from user  $X$ , it is also revoked from all users that were granted the privilege *solely* from  $X$

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## Trojan Horse Attack

- Suppose user *Bob* has privileges to read a secret table  $T$ .
  - User *Mallory* wants to see the data in  $T$  (but does not have the privileges to do so).
- 1 *Mallory* creates a table  $T'$  and gives INSERT privileges to *Bob*.
  - 2 *Mallory* tricks *Bob* into copying data from  $T$  to  $T'$  (e.g. by extending the “functionality” of a program used by *Bob*).
  - 3 *Mallory* can then see the data that comes from  $T$ .

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## Mandatory Access Control

### Idea

Achieve security based on system-wide policies that cannot be changed by individual users.

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## The Bell-LaPadula Model

- Basis: a partially ordered set of *security classes*
  - Example:  $TopSecret > Secret > Confidential > Unclassified$
- DB objects (e.g. tables, rows, columns) are assigned such a class
- Subjects (users, programs) are assigned *clearance* for such a class
- Goal: Information should never flow from a higher to a lower class.
- Restrictions enforced by the DBMS:
  - 1 Subject  $S$  can read object  $O$  only if  $clearance(S) \geq class(O)$
  - 2 Subject  $S$  can write object  $O$  only if  $clearance(S) \leq class(O)$

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## Trojan Horse Attack Revisited

- Suppose user *Bob* has privileges to read a secret table  $T$ .
  - $clearance(Bob) := Secret$
- User *Mallory* wants to see the data in  $T$  (but does not have the privileges to do so).
  - $clearance(Mallory) < Secret$
- 1 *Mallory* creates a table  $T'$  and gives INSERT privileges to *Bob*.
  - $class(T') := clearance(Mallory)$
  - i.e.  $class(T') < Secret$
- 2 *Mallory* tricks *Bob* into copying data from  $T$  to  $T'$ .
  - writing to  $T'$  **fails** for *Bob* because  $clearance(Bob) \not\leq class(T')$
- 3 *Mallory* can then see the data that comes from  $T$ .

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## Multilevel Relations

- Individual tuples or columns can be assigned security classes
  - ⇒ users with different clearances see different tables
- Example:  

ProjectEmployees			
EID	PID	EmpRole	Security Class
3	886	Manager	Unclassified
2	881	Researcher	TopSecret

  - Users with clearance TopSecret see two rows;
  - other users see only one.
- To avoid revealing **any** information, the Security Class attribute must be treated as part of the primary key.

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## Summary

- Three main security objectives:
  - Secrecy
  - Integrity
  - Availability
- Discretionary access control
  - based on notion of privileges
  - GRANT and REVOKE
  - susceptible to trojan horse attack
- Mandatory access control
  - based on notion of security classes
  - not widely supported

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