

Database Security

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Notes

Outline

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Granting and Revoking Privileges
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- 3 Mandatory Access Control
The Bell-LaPadula Model
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Related Concepts

Authentication: confirming the identity of users (or programs)

Authorization*: specifying access rights to resources

Encryption: encoding data to prevent unauthorized persons
from reading it (if they managed to access it)

*Our topic today.

Notes

Objectives in Securing a Database System

- Secrecy:** protection of data against unauthorized disclosure
- e.g. a student cannot see other students' grades

- Integrity:** prevention of unauthorized data modification
- e.g. only the instructor may assign grades

- Availability:** ensuring authorized access is possible
- e.g. students are not denied seeing their own grades

Notes

Access Control in a Database System

A **security policy** specifies who is authorized to do what in the system.

- A DBMS provides **access control** mechanisms to help implement a security policy.
- Two complementary types of mechanisms:
 - ① *Discretionary access control*
 - ② *Mandatory access control*

Notes

Discretionary Access Control

Idea
Achieve security based on:

- ① *privileges (certain access rights for tables, columns, etc.), and*
- ② *a mechanism for granting and revoking such privileges at a user's own discretion*

Authorization administration policy: specifies how granting/revoking of privileges is organized (i.e. who may grant and revoke)

- *Centralized administration:* only some privileged users
- *Ownership-based administration:* creator of the object

Administration delegation: If authorized to do so, a user may assign other users the right to grant or revoke.

In SQL-92, privileges are given to users.
In SQL:1999, privileges are given to *roles*; those are assigned to users.

Notes

Granting and Revoking Privileges in SQL

GRANT privileges ON object TO users [WITH GRANT OPTION]

- Possible privileges:
 - SELECT
 - INSERT (column)
 - UPDATE (column)
 - DELETE
 - REFERENCES (column)
- WITH GRANT OPTION allows user to pass on privilege (with or without passing on grant option)

REVOKE [GRANT OPTION FOR] privileges ON object
FROM users { RESTRICT | CASCADE }

- When a privilege is revoked from user X , it is also revoked from all users that were granted the privilege *solely* from X

Notes

Trojan Horse Attack

- Suppose user *Bob* has privileges to read a secret table T .
- User *Mallory* wants to see the data in T (but does not have the privileges to do so).

- 1 *Mallory* creates a table T' and gives INSERT privileges to *Bob*.
- 2 *Mallory* tricks *Bob* into copying data from T to T' (e.g. by extending the “functionality” of a program used by *Bob*).
- 3 *Mallory* can then see the data that comes from T .

Notes

Mandatory Access Control

Idea

Achieve security based on system-wide policies that cannot be changed by individual users.

Notes

The Bell-LaPadula Model

- Basis: a partially ordered set of *security classes*
 - Example: $TopSecret > Secret > Confidential > Unclassified$
- DB objects (e.g. tables, rows, columns) are assigned such a class
- Subjects (users, programs) are assigned *clearance* for such a class
- Goal: Information should never flow from a higher to a lower class.
- Restrictions enforced by the DBMS:
 - 1 Subject S can read object O only if $clearance(S) \geq class(O)$
 - 2 Subject S can write object O only if $clearance(S) \leq class(O)$

Notes

Trojan Horse Attack Revisited

- Suppose user *Bob* has privileges to read a secret table T .
 - $clearance(Bob) := Secret$
- User *Mallory* wants to see the data in T (but does not have the privileges to do so).
 - $clearance(Mallory) < Secret$
- 1 *Mallory* creates a table T' and gives INSERT privileges to *Bob*.
 - $class(T') := clearance(Mallory)$
 - i.e. $class(T') < Secret$
- 2 *Mallory* tricks *Bob* into copying data from T to T' .
 - writing to T' **fails** for *Bob* because $clearance(Bob) \not\leq class(T')$
- 3 *Mallory* can then see the data that comes from T .

Notes

Multilevel Relations

- Individual tuples or columns can be assigned security classes
 - ⇒ users with different clearances see different tables
- Example:

ProjectEmployees

EID	PID	EmpRole	Security Class
3	886	Manager	Unclassified
2	881	Researcher	TopSecret

 - Users with clearance TopSecret see two rows;
 - other users see only one.
- To avoid revealing **any** information, the Security Class attribute must be treated as part of the primary key.

Notes

Summary

- Three main security objectives:
 - Secrecy
 - Integrity
 - Availability
- Discretionary access control
 - based on notion of privileges
 - GRANT and REVOKE
 - susceptible to trojan horse attack
- Mandatory access control
 - based on notion of security classes
 - not widely supported

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